

Disjoint Sets in Graphs and its Application to Electrical Networks

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Abstract: Problem statement: One of the well known problems in Telecommunication and Electrical Power System is how to put Electrical Sensor Unit (ESU) in some various selected locations in the system. **Approach:** This problem was modeled as the vertex covering problems in graphs. The graph modeling of this problem as the minimum vertex covering set problem. **Results:** The degree covering set of a graph for every vertex is covered by the set minimum cardinality. The minimum of a graph cardinality of a degree covering set of a graph G is the degree covering number $\gamma^P(G)$. **Conclusion:** We show that Degree Covering Set (DCS) problem is NP-complete. In this study, we also give a linear algorithm to solve the DCS for trees. In addition, we investigate theoretical properties of $\gamma^P(T)$ in trees T .

Key words: Electrical Sensor Unit (ESU), vertex covering, Degree Covering Set (DCS), degree covering number, electrical networks, electrical power system, graph modeling

INTRODUCTION

At this moment we are dependable with the electricity power. The consumption of electricity is increasing every year. The electric power industries have to monitor carefully their system's state periodically as defined by a set of state variables (Ramos and Tahan, 2009; Kumkratug, 2010; Osuwa and Igwiro, 2010). One method of monitoring these variables is to place Electrical Sensor Unit (ESU) in some various selected locations in the system (Ahmadia *et al.*, 2010). Because of the high cost of a ESU, it is desirable to minimize their number while maintaining the ability to monitor (observe) the entire system (Ketabi and Hosseini, 2008). We define a system is to be observed if all of the state of the system can be decided from a set of measurements.

Approach: Let $G = (V, E)$ be a graph that modeled as an electrical system, where a vertex represents an electrical point (such as a electrical point of bus) and an edge represents a electrical line. The objective is how can we locate a smallest set of ESU to monitor the entire system is a graph theory problem closely related to the well-known minimum vertex covering problems. Hence, Electrical Sensor Unit (ESU) is not only of interest in the electrical engineering but also as a challenge problem in graph theory. For more discussion about covering and related subset problems as well as terminology not defined here, we encourage the reader to two books (Ramos and Tahan, 2009).

A ESU measures the state variable (voltage and phase angle) for the vertex at which it is placed and its

incident edges and their end vertices. The other finding rules are as follows:

- We observed the incidences of vertices to edges
- We observed joining any edges

MATERIALS AND METHOD

For a given set of vertices P representing the nodes where the ESU are placed, the following algorithm determines the collection of collection of vertices C and edges F .

Let $P \subseteq V$ be the set of vertices where the ESU are placed:

- Initialize $C = P$ and $F = \{e \in E(G) \mid e \text{ is incident to a vertex in } P\}$
- Add to C any vertex not already in C which is incident to an edge in F
- Add to F any edge not already in F such that
 - Both of its end vertices are in C or
 - It is incident to a vertex v of degree greater than one for which all the other edges incident to v are in F
- If steps 2 and 3 fail to locate any new edges or vertices for inclusion, stop. Otherwise, go to step 2

The objective of the problem is to solve the electrical system, we want $C = V(G)$ and $F = E(G)$ and we want to minimize $|P|$. This monitoring problem was introduced and studied in (Pask *et al.*, 2005; Katrenic and Semanisin, 2010).

A covering set in a graph $G = (V, E)$ is a set of $S \in V(G)$ if every vertex in $V - S$ has at least one neighbor in S . The cardinality of a minimum covering set of G is the covering number $\gamma(G)$.

Considering system monitoring their system's state as a modeling of the covering set problem, we define a set S to be a Degree Covering Set (DCS) if every vertex and every edge in G is observed by S . The Degree Covering Number $\gamma_P(G)$ is the minimum cardinality of a degree covering set of G . A covering set (respectively, degree covering set) of G with minimum cardinality is called a $\gamma(G)$ -set (respectively, $\gamma_P(G)$ -set).

We can conclude that covering set is a DCS, we have the following finding (Horaka and McAvaneyb, 2007).

Theorem 1: G is random graph and we conclude that, $1 \leq \gamma_P(G) \leq \gamma(G)$.

We will make next finding gives examples of graphs having a degree covering number equal to 1.

Theorem 2: G is random graph where $G \in \{K_n, C_n, P_n, K_{2,n}\}$ then $\gamma_P(G) = 1$.

The shadow of two random graphs G and H , denoted $G \circ H$, is the graph formed from one copy of G and $|V(G)|$ copies of H where the i -th vertex of G is adjacent to every vertex in the i -th copy of H . We can see that, the difference between $\gamma(G) - \gamma_P(G)$ can be large. For arbitrary Graph T , the shadow of a star $T = K_{1,k} \circ K_1$ has $\gamma_P(T) = 1 < k+1 = \gamma(T)$ for $k \geq 1$.

We note that every random graph H is the induced subgraph of a random graph G having $\gamma_P(G) = \gamma(G)$. Consider, for example, the shadow of random graph $G = (H \circ K_2)$, where $\gamma_P(G) = \gamma(G)$.

Hence, we have the following finding.

Theorem 3: Every random graph H is the induced subgraph of a random graph G having $\gamma_P(G) = \gamma(G)$.

Suppose that G is a random graph with maximum degree at least 3 and that S is a $\gamma_P(G)$ -set that contains a vertex v of degree less than 3. Let u be a vertex of degree at least 3 at minimum distance from v in random graph G , then, $(S - \{v\}) \cup \{u\}$ is also a minimum degree covering set of G . We can conclude as follows.

Theorem 4: If G is a random graph with maximum degree at least 3, then random graph G contains a $\gamma_P(G)$ -set that for every vertex has degree at least 3.

The vertices of small degree play a significant role in determining the degree covering number of a random graph. In particular, it is not necessarily true that if G' is the random graph obtained from a random graph G by subdividing one edge of G , then $\gamma_P(G') = \gamma_P(G)$.

In above Fig. 1, random tree T and T' are the trees that taken from T by dividing the edge vw once, then $\gamma_P(T') > \gamma_P(T)$.

In study (Sharaeh, 2008) Sharaeh investigated about the approximation algorithms for power system monitoring problem solution. We will investigate that the DCS problem is NP-complete even when restricted to bipartite or chordal graphs. We also investigate about linear time algorithm to find a DCS in trees and study theoretical properties of the degree covering number in trees.

NP-completeness: NP-Completeness is one the classic problem in Computer Science that find in many domain in electrical engineering. In this part we discuss the problem of NP-complete even when restricted to bipartite or chordal graphs (Horaka and McAvaneyb, 2007).

Degree Covering Set (Dcs):

Input: A graph $G = (V,E)$ and a positive integer $k > 1$.

Question: Does G have a DCS of size at most k ?

Base on the finding on (Kumkratug, 2010) we will make some modifications of the standard proof of NP-completeness of COVERING SET, that able to establish the NP-completeness of the DCS even when restricted to bipartite graphs.

Theorem 5: Degree covering set is NP-complete for bipartite graphs.

Proof: We first show that $DCS \in NP$. This is easy to do since one can verify a "yes" input of DCS in polynomial time. That is, for a random graph $G = (V,E)$, a positive integer k and for each subset $S \in V$ with $|S| \leq k$, it is easy to verify in polynomial time whether S is a DCS.

We next construct a reduction from NP-complete problem 3-DCS.

The problem of 3-DCS can be defines follows. Assume we have $V = \{v_1, v_2, \dots, v_n\}$ as set of variables and a set of 3-element sets $K = \{K_1, K_2, \dots, K_m\}$, called clauses, where each clause K_i contains three distinct occurrences of either a variable v_i or its complement \bar{v}_i .

We will solve whether K have a satisfying truth assignment, i.e., an assignment of TRUE and FALSE to the variables in V such that at least one variable (or its complement) in each clause $K_i \in K$ is assigned the value TRUE?

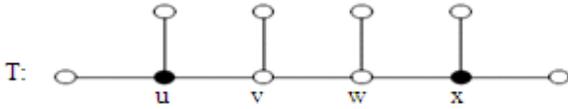


Fig. 1: Degree Covering Number of tree equal with 2

Given an input K of 3-DCS, we construct an input $G(K)$ of the DCS as follows. For each variable v_i , construct a cycle on four vertices K_4 , where two nonadjacent vertices are labeled v_i and v_i . These cycles and vertices are called random cycles and random vertices, respectively. For each clause $K_j = \{v_i, v_k, v_l\}$ create two nonadjacent vertices labeled $K_{j,1}$ and $K_{j,2}$ (called clause vertices) and add edges: $(v_i, K_{j,1})$, $(u_i, K_{j,2})$, $(v_k, K_{j,1})$, $(v_k, K_{j,2})$, $(v_l, K_{j,1})$, $(v_l, K_{j,2})$. By construction, the graph $G(K)$ is bipartite.

We will show that K has a satisfying truth assignment if and only if the graph $G(K)$ has a DCS of size at most $k = n$.

Assume that first that K has a satisfying truth assignment. We create a DCS P in $G(K)$ as follows: Assume that variable vertex v_i is map as TRUE and we map variable vertex v_i in P ; otherwise, put variable vertex v_i in P . The set P is a DCS for two reasons: (i) Base on the definition of DCS that v Vertex in each variable cycle K_4 belongs to P are dominated. Thus, every edge and vertex in each variable cycle K_4 is observed. (ii) Base on Theorem 5 that mention that for every clause vertex K_j is dominated by at least one vertex in P . We can conclude that, all edges between vertex, $K_{j,1}$ or $K_{j,2}$ and P is a DCS of size at most $k = n$.

Conversely, we must show that if $G(K)$ has a DCS of size at most $k = n$, then K has a satisfying truth assignment. Notice first that if P is a DCS of $G(K)$, then it must contain at least one vertex from each variable cycle K_4 . Therefore, $|P| \geq n$, i.e., $|P| = n$.

Theorem 6: Degree covering set is NP-complete for chordal graphs.

Tree: In this part, we investigate the degree covering number of a tree.

We will define some notations that we will use for the rest of the study.

For any vertex $v \in V$, the nearest neighborhood of v , denoted by $N(v)$, is the set $\{u \in V \mid uv \in E\}$ and its farthest neighborhood $N[v] = N(v) \cup \{v\}$. For a set $S \subseteq V$, its nearest neighborhood $N(S) = \cup_{v \in S} N(v)$ and its farthest neighborhood $N[S] = N(S) \cup S$. The personal neighbor set of a vertex v with respect to a set S , denoted $pn[v, S]$, is the set $N[v] - N[S - \{v\}]$. If $pn[v, S] \neq \emptyset$ for some vertex v and some $S \subseteq V$, then every vertex

of $pn[v, S]$ is called a personal neighbor of v with respect to S , or just an S -pn.

If T is a tree rooted at r and v is a vertex of T , then the level number of v , which we denote by $l(v)$, is the length of the unique r - v path in T . The definition of child and parent will be derived from concept of Tree. A vertex x is a descendant of y (and y is an ancestor of x) if the level lexicographic order of v - w path are monotonically increasing. We let $D(v)$ denote the set of descendants of v and we define $D[v] = D(v) \cup \{v\}$. The maximum of subtree of T rooted at u is the subtree of T induced by $D[u]$ and is denoted by T_u . We defined leaf as an end vertex of T . Downline vertex is a vertex adjacent to a leaf and a vertex adjacent to more than one leaf is called a Totally Downline Vertex.

Let T be the tree that divide any number of its edges of any number of times where T has at most one vertex of degree 2 or more. We define that tree as a shade tree.

Theorem 7: For any tree T , $\gamma_P(T) = 1$ if and only if T is a shade tree.

Proof: Suppose T is a shade tree. If T is a path, then any vertex of T forms a DCS in T . On the other hand, if T is not a path, then the vertex of maximum degree in T forms a DCS in T . In any event, $\gamma_P(T) = 1$. This proves the sufficiency.

To prove the necessity, suppose that T is not a shade tree. Then T contains at least two vertices, say u and v , of degree at least 3. We may assume that T is rooted at v . Let S be any DCS of T . If $|S| = 1$, then, renaming u and v if necessary, we may assume that no vertex in the maximum subtree T_u rooted at u belongs to S .

Base on this explanation, we conclude that, this is contradiction since no edge in T_u .

After this, we will characterize trees T with equal covering and degree covering numbers. We will use the following finding.

Theorem 8: If v is a totally downline vertex in a graph G , then v is in every $\gamma(G)$ -set and every $\gamma_P(G)$.

Proof: Thus assume that $\gamma(T) \geq 2$. Let T be a random tree with $\gamma(T)$ -set S , where every vertex in S is a totally downline vertex. Then Theorem 8 implies that S is also a unique $\gamma_P(T)$ -set. Hence, $\gamma_P(T) = \gamma(T)$.

For some $v \in S$ in $\gamma(T)$ -set S and that v is not a totally downline vertex.

With that $pn[v, S] = \{v\}$. Then every neighbor of v is covered by $S - \{v\}$ and v is an isolate in the random subgraph $\langle S \rangle$. Moreover, all vertices and edges are observed by $S - \{v\}$. Since T is a nontrivial tree, v has at least one neighbor, say u , in $V - S$. Then v is also

observed by $S-\{v\}$ and so $S-\{v\}$ is a DCS for T , contradicting the fact that $\gamma_P(T) = \gamma(T)$. We conclude, $|\text{pn}[v, S]-\{v\}| \geq 1$.

From this, we conclude that every vertex in S is a totally downline vertex and S is a unique covering set of T as claimed.

Before proceeding further, we define the shade tree number of a tree T , denoted by $\text{stn}(T)$, to be the minimum number of subsets into which $V(T)$ can be partitioned so that each subset induces a shade tree. We call such a partition a shade tree partition and each set of the partition a shade tree subset.

Theorem 9: For random tree T , $\text{stn}(T) \leq \gamma_P(T)$.

Proof: Suppose $m = 1$. Then, by Theorem 7, T is a shade tree and so $\text{stn}(T) = 1$. Suppose, then, that random trees T' with $\gamma_P(T') = m$, where $m \geq 1$, satisfy $\text{stn}(T') \leq \gamma_P(T')$. Let T be a tree with $\gamma_P(T) = m + 1$.

Let T be a random tree that has root at the vertex v_{m+1} . Our assumption is that v_1 is the vertex of S at maximum distance from v_{m+1} in tree T . Then u_1 has at least two neighbors and each neighbor of u_1 has degree at most 1 in tree T . Let u_1 be the ancestor of v_1 of degree at least 3 that is at a minimum distance from v_1 . It is not difficult to prove that u_1 is the neighbor of v_1 that has distance with degree 3 in the tree. Thus, $\text{stn}(T) \leq \gamma_P(T)$.

From above Theorem, we can conclude with this Theorem.

Theorem 10: For any tree T , $\gamma_P(T) = \text{stn}(T)$.

As a consequence of Theorem 12, we can determine a lower bound on the degree covering of a tree in terms of the number of vertices of degree at least 3.

Theorem 11: If T is a tree having k vertices of degree at:

$$\gamma_P(T) \geq \frac{k+2}{3}$$

Next we present a linear algorithm for finding a minimum DCS in a nontrivial tree T .

Algorithm FMDCS: Finding Minimum DCS on a non trivial tree

Input: A tree T on vertices more than 2 with the vertices labeled v_1, v_2, \dots, v_n so that $l(v_i) \leq l(v_j)$ for $i > j$.

Output: A minimum DCS S of G and a partition of $V(G)$ into $|S|$ subsets $\{V_x \mid x \in S\}$ so that each subset induces a shade tree.

Begin:

1. If G is a shade tree, then $S \leftarrow \{v_n\}$ and $V_{v_n} \leftarrow V(G)$ and output S and $\{V_x \mid x \in S\}$; otherwise, continue.
 2. If $\text{deg}_T v \leq 2$, then
 - 2.1. If there exists a child u of v (in G) such that $\text{Type}(u) = \text{TRUE}$ and $u \in V_x$ for some $x \in S$, then
 - 2.1.1. If v is a leaf (in T), then
 - 2.1.1.1. $V_x \leftarrow V_x \cup \{v\}$,
 - 2.1.1.2. $T \leftarrow T - v, I \leftarrow I \setminus \{i\}$ and go to step 10;
 - 2.1.2. if $\text{deg}_T v = 2$, then
 - 2.1.2.1. $V_x \leftarrow V_x \in V(T_v)$,
 - 2.1.2.2. $w \leftarrow v$ and go to step 12;
 - 2.2. Otherwise (if no such child exists), $I \leftarrow I \setminus \{i\}$ and go to step 10;
 3. $S \leftarrow S \cup \{v\}$.
 4. $w \leftarrow v_{m+1}$.
 5. $u \leftarrow \{\text{child of } w \text{ on the } w-v \text{ path}\}$.
 6. If $w = v_n$, then
 - 6.1. If the component of $T - uw$, then output S and $\{V_x \mid x \in S\}$;
 - 6.2. otherwise, $V_n \leftarrow V(T_u)$ and go to step 11.
 7. If $w \neq v_n$, then
 - 7.1. $z \leftarrow \text{parent}(w)$;
 - 7.2. If $\text{deg}_T w \geq 4$ or if $\text{deg}_T w = 3$ and the component of $T - \{uw, wz\}$ containing w is a not a path, then $V_v \leftarrow V(T_u)$ and go to step 11;
 8. $T \leftarrow T - V(T_u), I \leftarrow I \setminus \{k \mid v_k \in V(T_u)\}$ and go to step 13.
 9. $T \leftarrow T - V(T_w), I \leftarrow I \setminus \{k \mid v_k \in V(T_w)\}$, $\text{Type}(w) \leftarrow \text{FALSE}$. Go to step 7.
 10. $i \leftarrow \min\{k \mid k \in I\}$.
 11. If $i \leq n$, then return to step 4.
 12. If $i = n$, then
- End

We now verify the validity of Algorithm FMDCS.

Theorem 12: Algorithm FMDCS produces a $\gamma_P(G)$ -set in a nontrivial tree G .

Proof: Let G be a nontrivial tree of order n and let S be the set produced by Algorithm FMDCS. It follows from the way in which the set S is constructed that S is a DCS of G and so $\gamma_P(G) \leq |S|$. If G is a shade tree, then it follows from Theorem 7 and $S = \{v_n\}$ is a $\gamma_P(G)$ -set of G .

RESULTS

In this part we will test the scalability of the algorithm in Finding Minimum DCS in a nontrivial tree.

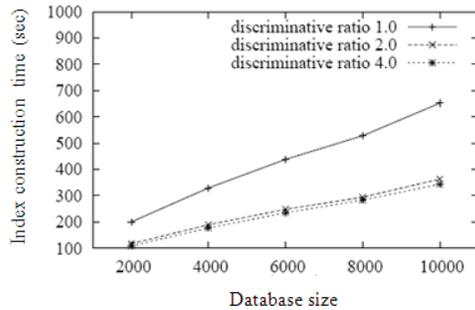


Fig. 2: The scalability of algorithm FMDCS

The scalability of algorithm FMDCS is presented in Fig. 2. We vary the database size from 2000-10000 and construct indexing as the number of minimum DCS from scratch for each database. We also define discriminative ratio thresholds to check how vary the tree inside the graph database.

DISCUSSION

We repeat the experiments for various minimum discriminative ratio thresholds. As shown in the figure, the index construction time is proportional to the database size.

Figure 2 also shows that the indexing construction time does not change too much for a given database when the minimum discriminative ratio is above 2.0. We find that the construction time consists of two parts, frequent graph mining and discriminative feature selection.

CONCLUSION

We consider the graph theoretical representation of this problem as a variation of the covering set problem and define a set S to be a degree covering set of a graph if every vertex and every edge in the system is monitored by the set S (following a set of rules for power system monitoring). The minimum cardinality of a degree covering set of a graph G is the degree covering number $\gamma_P(G)$. We show that the Degree Covering Set (DCS) problem is NP-complete even when restricted to bipartite graphs or chordal graphs.

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